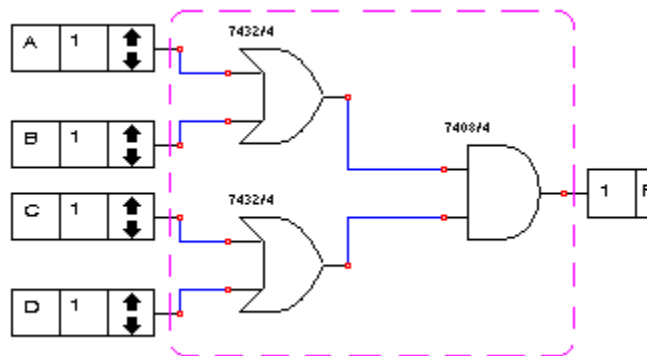


**CSCI 220: Computer Architecture-I**  
**Instructor: Pranava K. Jha**

**Design of Combinational Circuits**

**Q. 1.** Four large tanks at a chemical plant contain different liquids being heated. Liquid-level sensors are being used to detect whenever the level in tank *A* or tank *B* rises above a predetermined level. Temperature sensors in tanks *C* and *D* detect when the temperature in either of these tanks drops below a prescribed temperature limit. Assume that the liquid-level sensor outputs *A* and *B* are LOW when the level is satisfactory and HIGH when the level is too high. Also, the temperature-sensor outputs *C* and *D* are LOW when the temperature is satisfactory and HIGH when the temperature is too low. Design a logic circuit that will detect whenever the level in tank *A* or tank *B* is too high at the same time that the temperature in either tank *C* or tank *D* is too low.

It is not difficult to see that circuit output is given by  $(A + B)(C + D)$ .



**Q. 2.** Consider a four-input function  $F(A, B, C, D)$  that outputs a 1 whenever an odd number of its inputs are 1.

- Construct the truth table
- Generate the K-map. Can you simplify the function?
- Present an implementation of the function using two-input XOR gates.

Truth Table

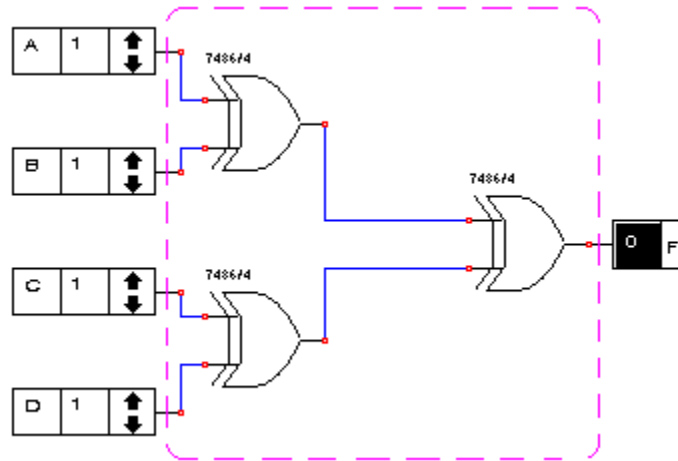
$A$	$B$	$C$	$D$	$F$
0	0	0	0	0
0	0	0	1	1
0	0	1	0	1
0	0	1	1	0
0	1	0	0	1
0	1	0	1	0
0	1	1	0	0
0	1	1	1	1
1	0	0	0	1
1	0	0	1	0
1	0	1	0	0
1	0	1	1	1
1	1	0	0	0
1	1	0	1	1
1	1	1	0	1
1	1	1	1	0

Karnaugh map is immediate.

	$C'D'$	$C'D$	$CD$	$CD'$
$A'B'$	0	1	0	1
$A'B$	1	0	1	0
$AB$	0	1	0	1
$AB'$	1	0	1	0

It is clear from the map that simplification is not possible. Further, it is not difficult to show that

$$F(A, B, C, D) = A \oplus B \oplus C \oplus D.$$



**Q. 3.** Construct a truth table for converting a four-bit binary string  $ABCD$  into a four-bit binary string  $WXYZ$  in *Gray code*. Use the scheme of building a sequence in Gray code that was given in the class. Further, obtain simplified expressions for  $W$ ,  $X$ ,  $Y$  and  $Z$  in terms of  $A$ ,  $B$ ,  $C$  and  $D$ .

Binary				Gray code			
$A$	$B$	$C$	$D$	$W$	$X$	$Y$	$Z$
0	0	0	0	0	0	0	0
0	0	0	1	0	0	0	1
0	0	1	0	0	0	1	1
0	0	1	1	0	0	1	0
0	1	0	0	0	1	1	0
0	1	0	1	0	1	1	1
0	1	1	0	0	1	0	1
0	1	1	1	0	1	0	0
1	0	0	0	1	1	0	0
1	0	0	1	1	1	0	1
1	0	1	0	1	1	1	1
1	0	1	1	1	1	1	0
1	1	0	0	1	0	1	0
1	1	0	1	1	0	1	1
1	1	1	0	1	0	0	1
1	1	1	1	1	0	0	0

It is clear that  $W = A$ . For  $X$ ,  $Y$  and  $Z$ , build Karnaugh maps to derive simplified expressions.

	$C'D'$	$C'D$	$CD$	$CD'$		$C'D'$	$C'D$	$CD$	$CD'$		$C'D'$	$C'D$	$CD$	$CD'$
$A'B'$	0	0	0	0	$A'B'$	0	0	1	1	$A'B'$	0	1	0	1
$A'B$	1	1	1	1	$A'B$	1	1	0	0	$A'B$	0	1	0	1
$AB$	0	0	0	0	$AB$	1	1	0	0	$AB$	0	1	0	1
$AB'$	1	1	1	1	$AB'$	0	0	1	1	$AB'$	0	1	0	1
<u>K-map for X</u>					<u>K-map for Y</u>					<u>K-map for Z</u>				

Simplified expressions:

$$\left| \begin{array}{l} W = A \\ X = A'B + AB' \\ \quad = A \oplus B \end{array} \right| \left| \begin{array}{l} Y = B'C + BC' \\ \quad = B \oplus C \end{array} \right| \left| \begin{array}{l} Z = C'D + CD' \\ \quad = C \oplus D \end{array} \right|$$

**Q. 4.** A four-variable logic function  $\alpha(A, B, C, D)$  is a 1 if any three or all four variables are equal to 1, and is a 0 otherwise. This is called the *majority function* of four variables. Obtain a minimum sum-of-products expression for  $\alpha(A, B, C, D)$ .

<u>Inputs</u>				<u>Output</u>
<i>A</i>	<i>B</i>	<i>C</i>	<i>D</i>	$\alpha$
0	0	0	0	0
0	0	0	1	0
0	0	1	0	0
0	0	1	1	0
0	1	0	0	0
0	1	0	1	0
0	1	1	0	0
0	1	1	1	1
1	0	0	0	0
1	0	0	1	0
1	0	1	0	0
1	0	1	1	1
1	1	0	0	0
1	1	0	1	1
1	1	1	0	1
1	1	1	1	1

Use a Karnaugh map to derive the following expression in minimum sum-of-products form:

$$\alpha(A, B, C, D) = ABC + ABD + ACD + BCD.$$

Here is the majority function of six variables  $A, B, C, D, E$  and  $F$ :

$$ABCD + ABCE + ABCF + ABDE + ABDF + ABEF + ACDE + ACDF + ACEF + ADEF + BCDE + BCDF + BCEF + BDEF + CDEF$$

**Q. 5.** Design a minimal two-level combinational network that detects the presence of any of the six illegal code groups in the 8421 code by providing a logic-1 output.

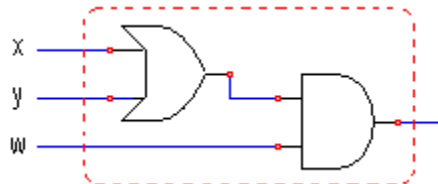
The truth table is immediate.

8	4	2	1	
$w$	$x$	$y$	$z$	$f$
0	0	0	0	0
0	0	0	1	0
0	0	1	0	0
0	0	1	1	0
0	1	0	0	0
0	1	0	1	0
0	1	1	0	0
0	1	1	1	0
1	0	0	0	0
1	0	0	1	0
1	0	1	0	1
1	0	1	1	1
1	1	0	0	1
1	1	0	1	1
1	1	1	0	1
1	1	1	1	1

The Karnaugh map follows.

	$y'z'$	$y'z$	$yz$	$yz'$
$w'x'$	0	0	0	0
$w'x$	0	0	0	0
$wx$	1	1	1	1
$wx'$	0	0	1	1

It is clear that  $f(w, x, y, z) = wx + wy = w(x + y)$ . A two-level implementation follows.



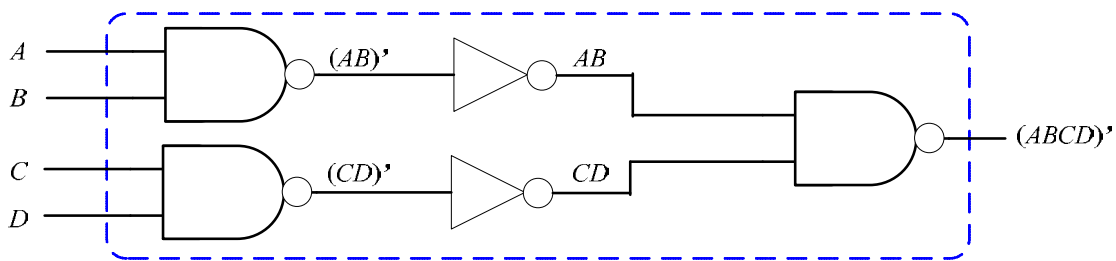
**Q. 6.** A 4-input NAND gate is defined by the following function:

$$f(A, B, C, D) = (ABCD)' = A' + B' + C' + D'$$

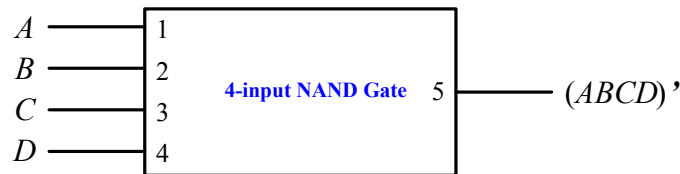
- Design a 4-input NAND gate using 2-input NAND gates and NOT gates. (Use as few gates as possible.)
- Design an 8-input NAND gate using two 4-input NAND gates and necessary 2-input NAND gates and NOT gates.

Note: IC-7420 houses dual 4-input NAND gates.

Here is an implementation of a 4-input NAND gate.



In order to build an 8-input NAND gate, let us first represent a 4-input NAND gate by means of the following block diagram.



An implementation of an 8-input NAND gate follows.

